



Bankruptcy Problems

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Summary

Bankruptcy problems

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Bankruptcy problems

Allocation of scarce resources (see O'Neill, 1982, Aumann and Maschler, 1985 and Curiel, Maschler and Tijs,1987)

$$BP = (N, c, E) = (E; c_1, ..., c_n)$$

 $\begin{array}{ll} \text{where} & N = \{1, ..., n\} & \text{set of claimants} \\ & c = (c_1, ..., c_n) \in \mathbb{R}^n_> & \text{vector of claims} \\ & E \in \mathbb{R}_\ge & \text{estate, with } E \leq \sum_{i \in N} c_i = C \end{array}$

A solution is an *n*-dimensional real vector $x = (x_1, ..., x_n)$, where x_i represents the monetary amount assigned to claimant $i \in N$, satisfying the following conditions:

$$0 \le x_i \le c_i, \qquad i \in N \qquad (rationality)$$

and

$$\sum_{i \in N} x_i = E \qquad \text{(efficiency)}$$

A solution rule is a function ψ that assigns a solution to each bankruptcy problem

Solutions

There exist four classical solutions (see Herrero and Villar, 2001)

• Proportional

PROP - The assignments are proportional to the claims:

$$PROP(N, c, E)_i = \frac{E}{C} c_i \qquad i \in N$$

• Constrained Equal Awards

CEA - The assignments are the same for all the agents, but not larger than the claims:

$$CEA(N, c, E)_i = \min\{\alpha, c_i\}$$
 $i \in N$

where α is such that $\sum_{i \in N} CEA(N, c, E)_i = E$

• Constrained Equal Losses

CEL - The assignments are equal to the claims reduced of the same amount for all the agents, but non-negative:

$$CEL(N, c, E)_i = \max\{c_i - \beta, 0\} \qquad i \in N$$

where β is such that $\sum_{i \in N} CEL(N, c, E)_i = E$

• Talmud

$$TAL(N, c, E) = \begin{cases} CEA(N, c/2, E) & \text{if } E \le C/2 \\ c/2 + CEL(N, c/2, E - C/2) & \text{if } E > C/2 \end{cases}$$

Example 1 (Solutions) Given the bankruptcy problem (72; 6, 9, 24, 33, 36)

$$C = 108; \frac{E}{C} = \frac{2}{3}$$

$$PROP = (4, 6, 16, 22, 24)$$

$$CEA = (6, 9, 19, 19, 19) \qquad [\alpha = 19]$$

$$CEL = (0, 1.5, 16.5, 25.5, 28.5) \qquad [\beta = 7.5]$$

$$TAL = (3, 4.5, 14.5, 23.5, 26.5)$$

PROP is the most intuitive CEA favors smaller claims CEL favors larger claims TAL is related to the nucleolus of the bankruptcy game

Duality

Given a solution rule ψ , the dual solution rule ψ^* produces the same solution when used for allocating the losses w.r.t. the claims:

$$\psi^*(N,c,E) = c - \psi(N,c,C-E)$$

CEA and CEL are dual rules, while PROP and TAL are *self-dual*:

- CEA(N, c, E) = c CEL(N, c, C E)
- CEL(N, c, E) = c CEA(N, c, C E)
- PROP(N, c, E) = c PROP(N, c, C E)
- $\bullet \; TAL(N,c,E) = c TAL(N,c,C-E)$

Comments

Very simple model of real-world situations

Most important additional elements in the existing literature:

- different priorities of the claimants (Young, 1994, Bebchuck and Fried, 1996, Schwarcz, 1997 and Kaminski, 2000)
- minimal rights of the claimants (Curiel, Maschler and Tijs, 1987 and Pulido, Sanchez-Soriano and Llorca, 2002)
- multiple issues (Calleja, Borm, Hendrickx, 2005 and Moreno-Ternero, 2009)
- negative claims and estate (Herrero, Maschler and Villar, 1999, Branzei, Ferrari, Fragnelli and Tijs, 2008 and 2011)
- non-transferable utility situations (Orshan, Valenciano and Zarzuelo, 2003 and Carpente, Casas-Mendez, Gozalvez, Llorca, Pulido and Sanchez-Soriano, 2013)

Two important surveys are due to Thomson (2003 and 2015)

Bankruptcy Rules and Min Cost Flow Problems Standard flow problem

Let G(N, A) be a network with two particular nodes, the source s with no entering arcs and the sink t with no outgoing arcs; arcs have minimal and maximal capacity constraints



A flow is a function $x: A \to \mathbb{R}_+$ that respects the capacity constraints and such that

$$\sum_{j \in N} x_{ij} = \sum_{j \in N} x_{ji}, \forall i \in N \setminus \{s, t\}$$

A classical bankruptcy problem can be represented as a standard flow problem



Each feasible flow corresponds to a solution of the bankruptcy problem

Min cost flow approach with suitable cost functions $k_i, i \in N$ can lead to classical rules, or suggest new divisions according to different fairness criteria, tailoring the amount on each agent

- PROP: $k_i(x_i) = \frac{x_i^2}{c_i}, i \in N$
- CEA: $k_i(x_i) = x_i^2, i \in N$
- CEL: $k_i(x_i) = (c_i x_i)^2, \ i \in N$
- TAL:
 - $\text{ if } E \leq \frac{1}{2} \sum_{i \in N} c_i$:

$$k_i(x_i) = x_i^2, \ i \in N$$

setting the maximal capacity of the arcs corresponding to the claimants to $\frac{1}{2}c_i$

- if
$$E > \frac{1}{2} \sum_{i \in N} c_i$$
:
 $k_i(x_i) = (c_i - x_i)^2, \ i \in N$

setting the minimal capacity of the arcs corresponding to the claimants to $\frac{1}{2}c_i$

Defining the minimal right of agent $i \in N$ as

$$m_i(N, c, E) = \max\{0, E - \sum_{j \in N \setminus \{i\}} c_j\}, \ i \in N$$

it is possible to introduce the *Adjusted Proportional* rule (Curiel, Maschler and Tijs, 1987):

$$APROP_i(N, c, E) = m_i(N, c, E) + PROP_i(Nc', E'), \ i \in N$$

where $E' = E - \sum_{j \in N} m_j$ and $c'_i = min\{E', c_i - m_i(N, c, E)\}, i \in N$ It coincides with the τ -value (Tijs, 1981) of the bankruptcy game

The corresponding cost functions are:

$$k_i(x_i) = \begin{cases} 0 & \text{if } x_i \le m_i \\ \frac{(x_i - m_i)^2}{c'_i - m_i} & \text{if } x_i > m_i \end{cases}, \ i \in N$$

Example 2 Consider the bankruptcy problem with $N = \{1, 2\}; c = (4, 12); E = 10$. The flow problems associated to the above five rules are depicted below (the notations for the arcs are "min capacity/max capacity" above the arc and "cost function" below the arc)



$$k_2(x_2) = \begin{cases} 0 & \text{if } x_2 \le 6\\ \frac{(x_2 - 6)^2}{4} & \text{if } x_2 > 6 \end{cases}$$

Bankruptcy Rules and Hydraulic Systems

According to Kaminski (2000) "hydraulic" rules can be represented as a system of connected vessels

• PROP rule



 \bullet *CEA* rule



• CEL rule



• TAL rule



• APROP rule



Bankruptcy Games

It is possible to define two TU-games, the pessimistic, (N, v_P) , and the optimistic, (N, v_O) :

$$v_P(S) = max\left(0, E - \sum_{i \in N \setminus S} c_i\right) \qquad S \subseteq N$$
$$v_O(S) = min\left(E, \sum_{i \in S} c_i\right) \qquad S \subseteq N$$

Example 3 (Inconsistency of the optimistic game) Given the bankruptcy problem (5; 3, 4), the two games are:

$$v_O(1) = 3; v_O(2) = 4; v_O(12) = 5$$

 $v_P(1) = 1; v_P(2) = 2; v_P(12) = 5$

The optimistic game assigns to the players as singletons 3 and 4, even if the estate is 5

The core of (N, v_P) coincides with the set of rational solutions of the bankruptcy problem:

$$x \in core(v_P) \iff \begin{cases} a) & \sum_{i \in N} x_i = E \\ b) & 0 \le x_i \le c_i, \quad i \in N \end{cases}$$

Game theoretic bankruptcy rules

A bankruptcy rule ψ is a game theoretic rule if it is possible to construct a solution concept F for cooperative games such that

 $\psi(N, c, E) = F(N, v_P)$

for all bankruptcy problems (N,c,E), where (N,v_P) is the pessimistic TU-game associated to (N,c,E)

Curiel, Maschler and Tijs (1987) proved that a bankruptcy rule ψ is a game theoretic rule if and only if *truncation property* holds, i.e.

$$\psi(N,c,E) = \psi(N,\bar{c},E)$$

where $\bar{c}_i = min\{c_i, E\}, i \in N$